

# Block I/O Layer Tracing: *blktrace*

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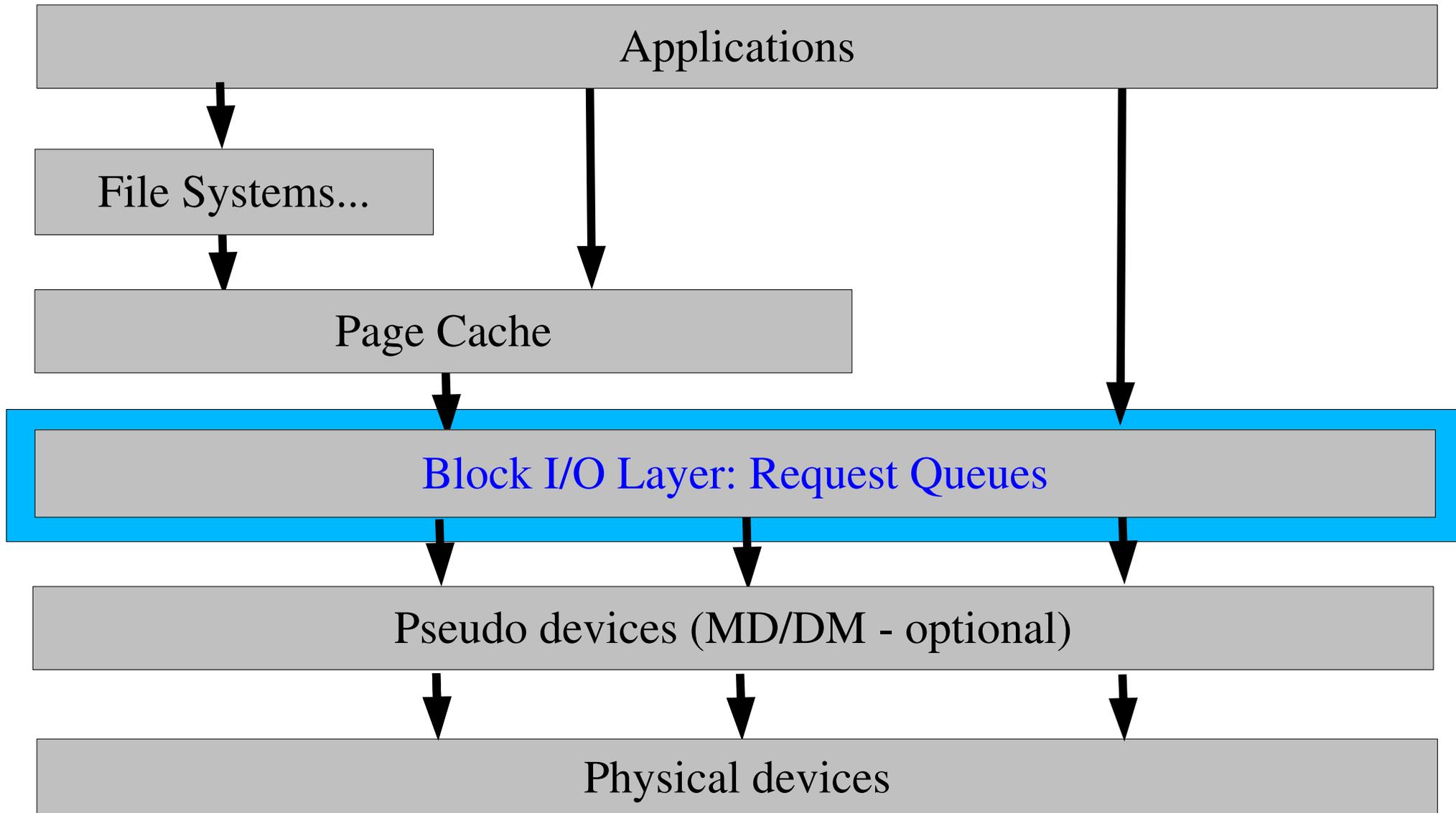
# Introduction

- Blktrace – overview of a new Linux capability
  - Ability to see what's going on inside the block I/O layer
    - “You can't count what you can't measure”
  - Kernel implementation
  - Description of user applications
- Sample Output & Analysis

# Problem Statement

- Need to know the specific operations performed upon *each* I/O submitted into the block I/O layer
- Who?
  - Kernel developers in the I/O path:
    - Block I/O layer, I/O scheduler, software RAID, file system, ...
  - Performance analysis engineers – *HP OSLO S&P*...

# Block I/O Layer (simplified)



# iostat

- The iostat utility *does* provide information pertaining to request queues associated with specific devices
  - Average I/O time on queue, number of merges, number of blocks read/written, ...
- However, it does *not* provide detailed information on a per-I/O basis

# Blktrace – to the rescue!

- Developed and maintained by Jens Axboe (block I/O layer maintainer)
  - My additions included adding threads & utility splitting, DM remap events, blkrawverify utility, binary dump feature, testing, kernel/utility patches, and documentation.
- Low-overhead, configurable kernel component which emits *events* for specific operations performed on each I/O entering the block I/O layer
- Set of tools which extract and format these events

However, *blktrace* is **not** an analysis tool!

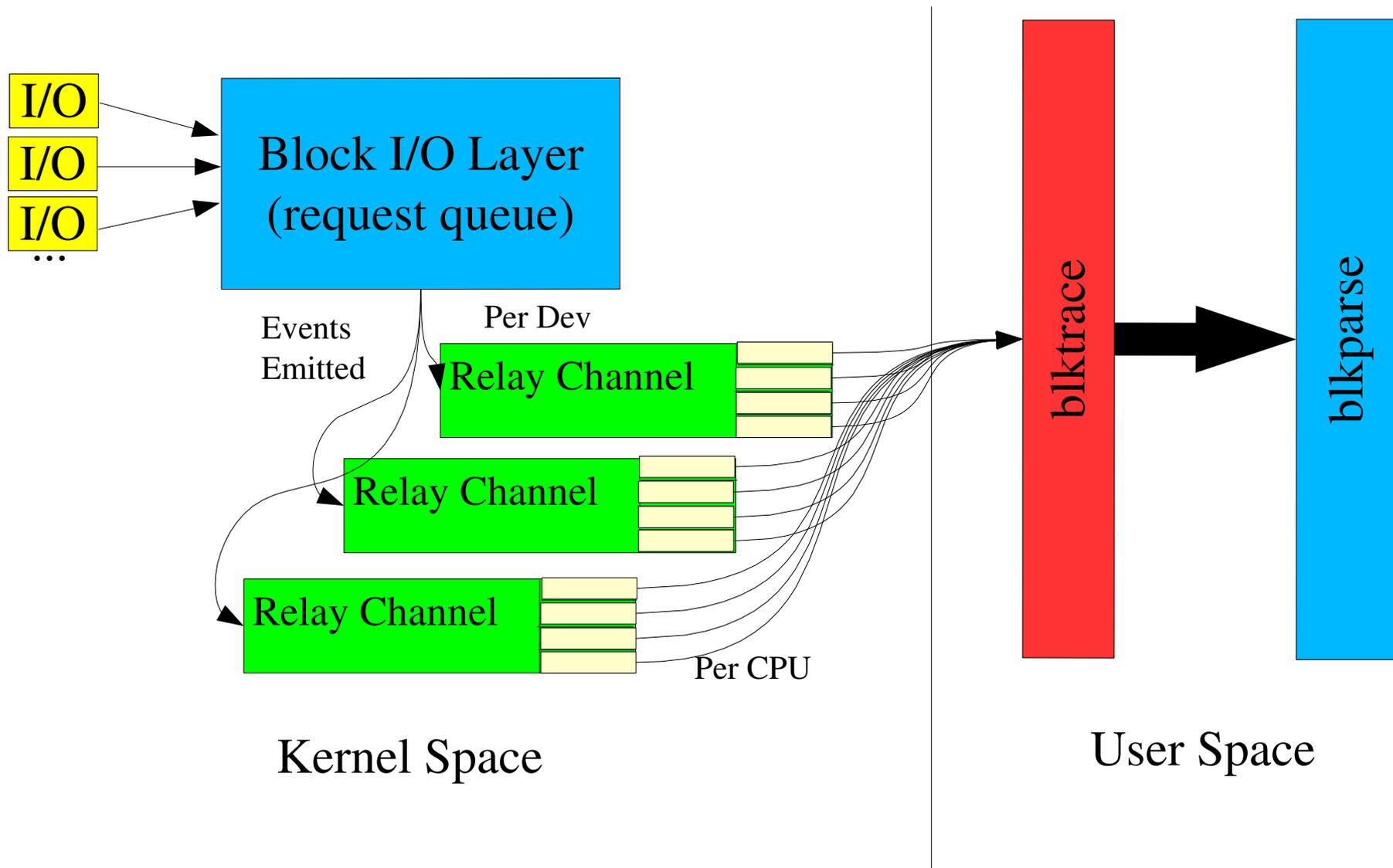
# Feature List

- Provides detailed block layer information concerning individual I/Os
- Low-overhead kernel tracing mechanism
  - Seeing less than 2% hits to application performance in relatively stressful I/O situations
- Configurable:
  - Specify 1 or more physical or logical devices (including MD and DM (LVM2))
  - User-selectable events – can specify filter at event acquisition and/or when formatting output
- Supports both “live” and “playback” tracing

# Events Captured

- Request queue entry allocated
- Sleep during request queue allocation
- Request queue insertion
- Front/back merge of I/O on request queue
- Re-queue of a request
- Request issued to underlying block dev
- Request queue plug/unplug op
- I/O split/bounce operation
- I/O remap
  - MD or DM
- Request completed

# *blktrace*: General Architecture



# *blktrace* Utilities

- *blktrace*: Device configuration, and event extraction utility
  - Store events in (long) term storage
  - Or, pipe to *blkparse* utility for live tracing
    - Also: networking feature to remote events for parsing on another machine
- *blkparse*: Event formatting utility
  - Supports textual or binary dump output

# *blktrace*: Event Output

```
Dev <mjr, mnr>
% blktrace -d /dev/sda -o - | blkparse -i -
  8,0 3 1 0.0000000000 697 G W 223490 + 8 [kjournald]
  8,0 3 2 0.000001829 697 P R [kjournald]
  8,0 3 3 0.000002197 697 Q W 223490 + 8 [kjournald]
  8,0 3 4 0.000005533 697 M W 223498 + 8 [kjournald]
CPU 8,0 3 5 0.000008607 697 M W 223506 + 8 [kjournald]
...
  8,0 3 10 0.000024062 697 D W 223490 + 56 [kjournald]
  8,0 1 11 0.009507758 0 C W 223490 + 56 [0]
```

Process

Sequence Number

Time Stamp

PID

Event

Start block + number of blocks

CPU

# blktrace: Summary Output

Per CPU details

Writes submitted on

CPU0 (sdao):  
Reads Queued: 0, 0KiB Writes Queued: 77,382, 5,865MiB  
Read Dispatches: 0, 0KiB Write Dispatches: 7,329, 3,020MiB  
Reads Requeued: 0 Writes Requeued: 6  
Reads Completed: 0, 0KiB Writes Completed: 0, 0KiB  
Read Merges: 0 Write Merges: 68,844  
Read depth: 2 Write depth: 65  
IO unplugs: 414 Timer unplugs: 414

...  
CPU3 (sdao):  
Reads Queued: 105, 18KiB Writes Queued: 14,541, 2,578MiB  
Read Dispatches: 22, 60KiB Write Dispatches: 6,207, 1,964MiB  
Reads Requeued: 0 Writes Requeued: 1,408  
Reads Completed: 22, 60KiB Writes Completed: 12,300, 5,059MiB  
Read Merges: 83 Write Merges: 10,968  
Read depth: 2 Write depth: 65  
IO unplugs: 287 Timer unplugs: 287

Total (sdao):  
Reads Queued: 105, 18KiB Writes Queued: 92,546, 8,579MiB  
Read Dispatches: 22, 60KiB Write Dispatches: 13,714, 5,059MiB  
Reads Requeued: 0 Writes Requeued: 1,414  
Reads Completed: 22, 60KiB Writes Completed: 12,300, 5,059MiB  
Read Merges: 83 Write Merges: 80,246  
IO unplugs: 718 Timer unplugs: 718

Throughput (R/W): 0KiB/s / 39,806KiB/s  
Events (sdao): 324,011 entries  
Skips: 0 forward (0 - 0.0%)

Per device details

Writes completed on

Avg throughput

# *blktrace*: Event Storage Choices

- Physical disk backed file system
  - *Pros*: large/permanent amount of storage available; supported by all kernels
  - *Cons*: potentially higher system impact; may negatively impact devices being watched (if storing on the same bus that other devices are being watched on...)
- RAM disk backed file system
  - *Pros*: predictable system impact (RAM allocated at boot); less impact to I/O subsystem
  - *Cons*: limited/temporary storage size; removes RAM from system (even when not tracing); may require reboot/kernel build
- TMPFS
  - *Pros*: less impact to I/O subsystem; included in most kernels; only utilizes system RAM while events are stored
  - *Cons*: limited/temporary storage; impacts system predictability – RAM “removed” as events are stored – could affect application being “watched”

## *blktrace: Analysis Aid*

- As noted previously, *blktrace* does not analyze the data; it is responsible for storing and formatting events
- Need to develop post-processing analysis tools
  - Can work on formatted output or binary data stored by *blktrace* itself
  - Example: *btt* – block trace timeline

# Practical *blktrace*

- Here at HP OSLO S&P, we are investigating I/O scalability at various levels
  - Including the efficiency of various hardware configurations and the effects on I/O performance caused by software RAID (MD and DM)
- *blktrace* enables us to determine scalability issues within the block I/O layer and the overhead costs induced when utilizing software RAID

# Life of an I/O *(simplified)*

- I/O enters block layer – it can be:
  - Remapped onto another device (MD, DM)
  - Split into 2 separate I/Os (alignment, size, ...)
  - Added to the request queue
  - Merged with a previous entry on the queue

*All I/Os end up on a request queue at some point*

- At some later time, the I/O is issued to a device driver, and submitted to a device
- Later, the I/O is completed by the device, and its driver

# *btt*: Life of an I/O

- Q2I – time it takes to process an I/O *prior* to it being inserted or merged onto a request queue
  - Includes split, and remap time
- I2D – time the I/O is “idle” on the request queue
- D2C – time the I/O is “active” in the driver and on the device
- $Q2I + I2D + D2C = Q2C$ 
  - Q2C: Total processing time of the I/O

# btt: Partial Output

Merge Ratio:  
#Q / #D

"Software" time

DEV	#Q	#D	Ratio	BLKmin	BLKavg	BLKmax	Total
[ 8, 0]	92827	12401	7.5	1	109	1024	10120441
[ 8, 1]	93390	13676	6.8	1	108	1024	10150343
[ 8, 2]	92366	13052	7.1	1	109	1024	10119302
[ 8, 3]	92278	13616	6.8	1	109	1024	10119043
[ 8, 4]	92651	13736	6.7	1	109	1024	10119903

DEV	Q2I	I2D	D2C	Q2C
[ 8, 0]	0.049697430	0.302734498	0.074038617	0.400079555
[ 8, 1]	0.031665593	0.050032148	0.058669682	0.125934697
[ 8, 2]	0.035651772	0.031035436	0.047311659	0.096735504
[ 8, 3]	0.021047776	0.011161007	0.038519804	0.060975408
[ 8, 4]	0.028985217	0.008397228	0.034344640	0.058160497

Avg I/O time

SCSI bus, target:  
low- to high-priority

DEV	Q2I	I2D	D2C
[ 8, 0]	11.7%	71.0%	17.4%
[ 8, 1]	22.6%	35.6%	41.8%
[ 8, 2]	31.3%	27.2%	41.5%
[ 8, 3]	29.8%	15.8%	54.5%
[ 8, 4]	40.4%	11.7%	47.9%

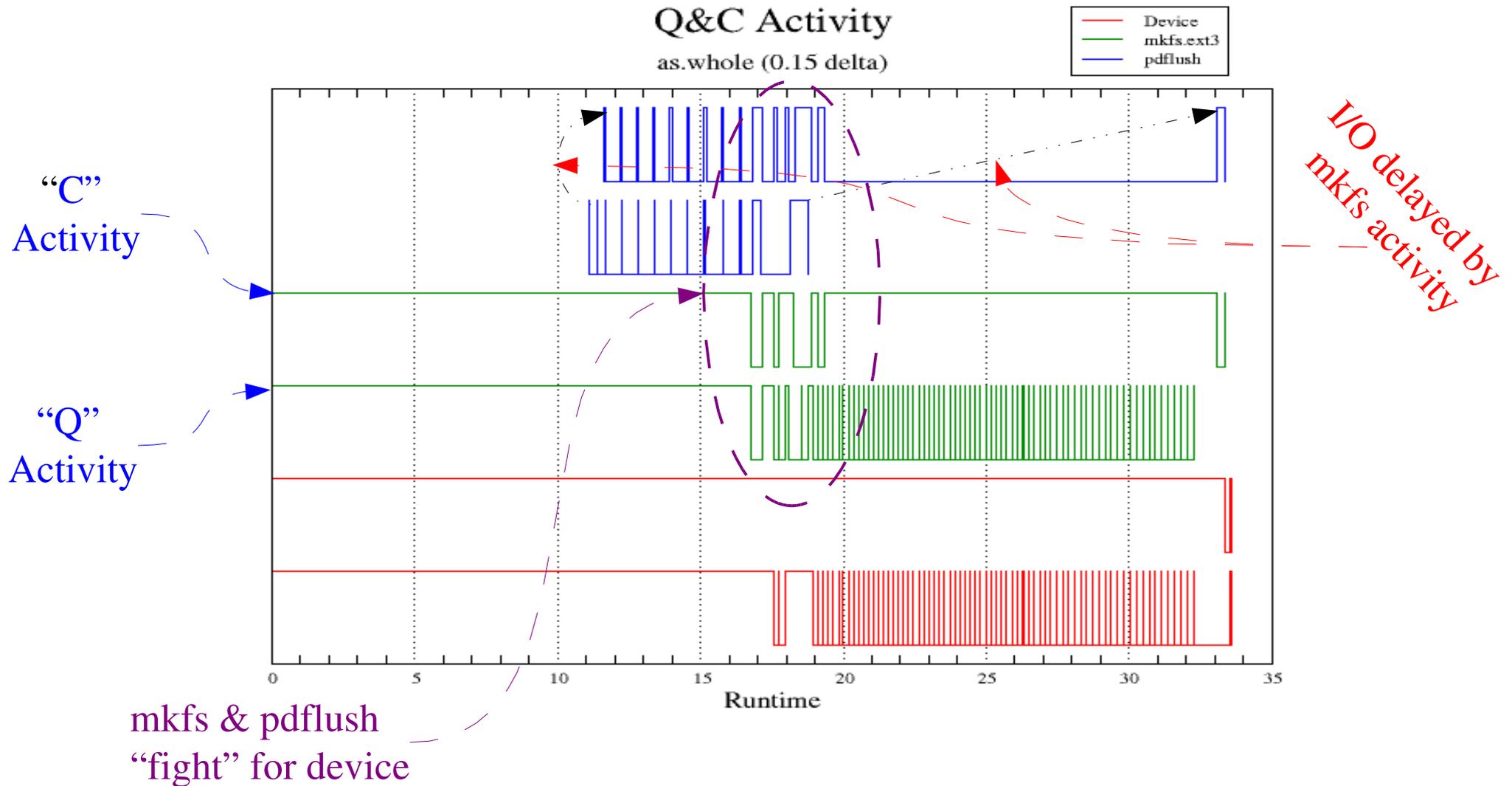
Excessive "idle" time on  
request queue

Driver/Device time

## *btt*: Q&C Activity

- *btt* also generates “activity” data – indicating ranges where processes and devices are actively handling various events (block I/O entered, I/O inserted/merged, I/O issued, I/O complete, ...)
- This data can be plotted (e.g. xmgrace) to see patterns and extract information concerning anomalous behavior

# *btt*: I/O Scheduler Example



# btt: I/O Scheduler - Explained

- Characterizing I/O stack
- Noticed very long I2D times for certain processes
- Graph shows continuous stream of I/Os...
  - ...at the device level
  - ...for the *mkfs.ext3* process
- Graph shows significant lag for *pdflush* daemon
  - Last I/O enters block I/O layer around 19 seconds
  - But: last batch of I/Os aren't completed until 14 seconds later!
- Cause? Anticipatory scheduler – allows *mkfs.ext3* to proceed, holding off *pdflush* I/Os

# Resources

- Kernel sources:
  - Patch for Linux 2.6.14-rc3 (or later, up to 2.6.17)
  - Linux 2.6.17 (or later) – built in
- Utilities & documentation (& kernel patches)
  - `rsync://rsync.kernel.org/pub/scm/linux/kernel/git/axboe/blktrace.git`
  - See documentation in *doc* directory
- Mailing list: [linux-btrace@vger.kernel.org](mailto:linux-btrace@vger.kernel.org)